## **TRALE Definite Clauses** Grammar Engineering, SS 2006 Georgiana Dinu





### Overview Definite clauses: Overview Clauses in Prolog Language TRALE Definite Clauses Defining a clause Relational Attachments

```
Definite Clauses
   • (1) \phi_1 \wedge \ldots \wedge \phi_n \rightarrow \phi_0
            \phi_0 \quad :-\phi_1, \dots, \phi_n
            head: -body
      (2) \phi_0
    TRALE syntax
        <clause> ::= <literal> if <goal>.
        <literal> ::= <pred_sym>
                        <pred_sym>(<seq(desc)>)
      member2 as a definite clause in TRALE:
        member(X, hd:X) if true.
        member(X, tl:Xs) if member(X, Xs).
```

```
Definite Clauses in Prolog: Declarative
Semantics
     concatenate([], L, L).
     concatenate([X|L1], L2, [X|L3]) :- concatenate(L1, L2, L3).
     ?- concatenate([a], [b], [a,b]).
     concatenate([a], [b], [a,b]) :- concatenate([], [b], [b]).
   ^{ullet} A goal is true if it is the head of some clause instance and each of the goals (if
```

```
Definite Clauses in Prolog: Declarative
Semantics
     concatenate([], L, L).
     concatenate([X|L1], L2, [X|L3]) :- concatenate(L1, L2, L3).
     ?- concatenate([a], [b], [a,b]).
     concatenate([a], [b], [a,b]) :- concatenate([], [b], [b]).
     A goal is true if it is the head of some clause instance and each of the goals (if
```

any) in the body of that clause instance is true, where an instance of a clause (or term) is obtained by substituting, for each of zero or more of its variables, a new term for all occurrences of the variable.

Definite Clauses in Prolog: Procedural
Semantics
Query
?- concatenate(X, Y, [a,b]).
Instantiated queries
<pre>?- concatenate(X1, Y, [b]). ?- concatenate(X2, Y, []).</pre>
Solutions
X = [a,b] Y = []
X = [a] Y = [b]
X = []  Y = [a,b]
If there is no matching head for a goal, the execution backtracks to the most recent successful match.

. . . . . . . .

• •

#### The Cut Symbol

 The cut operation commits the system to all choices made since the parent goal was invoked, and causes other alternatives to be discarded.

```
member(X, [X|_]).
  member(X, [\_|L]) := member(X, L).
   ?- member(X, [d,e,f]).
  Solution: d, e, f.
  member(X, [X]) :- !.
  member(X, [\_|L]) :- member(X, L).
  ?- member(X, [d,e,f]).
  Solution: d
?- member(e, [d, e, f]) .
  Yes
```

```
If then else
    • P -> O ; R
      Analogous to if P then Q else R and defined as if by
      (P -> Q; R) :- P, !, Q.
      (P -> Q; R) :- R.
     Only explores the first solution to the goal P (removes all choice-points created by
      P and executes Q.)
    • P -> O
      When occurring as a goal, this construction is read as equivalent to:
      (P -> Q; fail)
      Example
      add(X,L1,L2) :- member(X,L1) -> L2 = L1 ; L2 = [X|L1].
       1 ?- add(a, [b,c], [a,b,c]) .
       Yes
       2 ?- add(X, [b,c], [a,b,c]) .
       NO
```

#### Negation by failure

 Fails if the goal P has a solution, and succeeds otherwise. This is not real negation ("P is false"), but a kind of pseudo-negation meaning "P is not provable". It is defined as:

```
\+(P) :- P, !, fail.
\+(_).
```

The only time we get something like the desired result if there is no existentially quantified variable in the goal.

```
Example
```

```
p(a).
p(b).
```

```
P(\mathbf{z})
```

```
q(c).
```

```
(-q(X), +p(X)). % succeeds with X=c
(-+p(X), q(X)). % fails
```

TRALE Definite Clauses
Terms: TRALE Descriptions
Evaluation of a definite clause query
Solution to Queries: Satisfiers of the entire query as most general solutions
<ul> <li>Example</li> </ul>
<pre>?- query member(X, [noun, verb]).</pre>
Solution 1:
member([0] noun
CASE case,
ne_list
HD [0]
TL ne_list
HD verb
VFORM vform
TL e_list)

• Terms: <sup>-</sup>	TRALE Descriptions
Evaluati	on of a definite clause query
Sol	ution to Queries: Satisfiers of the entire query as most general solutions
Exa	ample
	<pre>?- query member(X, [noun, verb]).</pre>
Sol	ution 2:
n	nember([0] verb
	VFORM vform,
	ne_list
	HD noun
	CASE case
	TL ne_list
	HD [0]
	TL e_list)

#### TRALE Definite Clauses: Syntax <goal> ::= true <literal> | prolog(<prolog\_goal>) | (<goal>,<goal>) | (<goal>;<goal>) | (<desc> =@ <desc>) | (<cut\_free\_goal> -> <goal>) | (<cut\_free\_goal> -> <goal> ; <goal>) | (\+ <goal>) | when(cond, <goal>)

```
TRALE Definite Clauses: Syntax
   <goal> ::= true
         <literal>
         | prolog(<prolog_goal>)
        | (<goal>,<goal>)
        | (<goal>;<goal>)
        | (<desc> =@ <desc>)
        | (<cut_free_goal> -> <goal>)
        | (<cut_free_goal> -> <goal> ; <goal>)
        | (\+ <goal>)
         | when(cond, <goal>)
```



```
TRALE Definite Clauses: Syntax
   <goal> ::= true
         l <literal>
         | prolog(<prolog_goal>)
        | (<goal>,<goal>)
        | (<goal>;<goal>)
        | (<desc> =@ <desc>)
        | (<cut_free_goal> -> <goal>)
        | (<cut_free_goal> -> <goal> ; <goal>)
        | (\+ <goal>)
        | when(cond, <goal>)
```

```
TRALE Definite Clauses: Syntax
   <goal> ::= true
         l <literal>
         | prolog(<prolog_goal>)
        | (<goal>,<goal>)
         | (<goal>;<goal>)
        | (<desc> =@ <desc>)
        | (<cut_free_goal> -> <goal>)
        | (<cut_free_goal> -> <goal> ; <goal>)
        | (\+ <goal>)
        | when(cond, <goal>)
```

```
TRALE Definite Clauses: Syntax
   <goal> ::= true
        l <literal>
        | prolog(<prolog_goal>)
        | (<goal>,<goal>)
        | (<goal>;<goal>)
        | (<desc> =@ <desc>)
        | (<cut_free_goal> -> <goal>)
        | (<cut_free_goal> -> <goal> ; <goal>)
        | (\+ <goal>)
        | when(cond, <goal>)
```

```
TRALE Definite Clauses: Syntax
   <goal> ::= true
        l <literal>
        | prolog(<prolog_goal>)
        | (<goal>,<goal>)
        | (<goal>;<goal>)
        (<desc> =@ <desc>)
        (<cut_free_goal> -> <goal>)
        | (<cut_free_goal> -> <goal> ; <goal>)
        | (\+ <goal>)
        | when(cond, <goal>)
```

```
TRALE Definite Clauses: Syntax
   <goal> ::= true
        | <literal>
        | prolog(<prolog_goal>)
        | (<goal>,<goal>)
        | (<goal>;<goal>)
        | (<desc> =@ <desc>)
        | (<cut_free_goal> -> <goal>)
        (<cut_free_goal> -> <goal> ; <goal>)
        | (\+ <goal>)
        | when(cond, <goal>)
```

#### TRALE Definite Clauses: Syntax <goal> ::= true | <literal> | prolog(<prolog\_goal>) | (<goal>,<goal>) | (<goal>;<goal>) | (<desc> =@ <desc>) | (<cut\_free\_goal> -> <goal>) | (<cut\_free\_goal> -> <goal> ; <goal>) | (\+ <goal>) | when(cond, <goal>)

#### TRALE Definite Clauses: Syntax <goal> ::= true | <literal> | prolog(<prolog\_goal>) | (<goal>,<goal>) | (<goal>;<goal>) | (<desc> =@ <desc>) | (<cut\_free\_goal> -> <goal>) | (<cut\_free\_goal> -> <goal> ; <goal>) ( + < goal >)when(cond, <goal>)

#### TRALE Definite Clauses: Syntax <goal> ::= true | <literal> | prolog(<prolog\_goal>) | (<goal>,<goal>) | (<goal>;<goal>) | (<desc> =@ <desc>) | (<cut\_free\_goal> -> <goal>) | (<cut\_free\_goal> -> <goal> ; <goal>) | (\+ <goal>) when(cond, <goal>)

```
Co-routining 1
    Syntax
             when(<cond>,<goal>)
    when/2 delays execution until some event is witnessed.
     Example:
     append(X,Y,Z) if
        when( ( X=(e_list;ne_list)
                ; Y=e_list
                ; Z=(e list;ne list)
                 undelayed_append(X,Y,Z)).
     undelayed_append(L,[],L) if true.
     undelayed_append([],(L,ne_list),L) if true.
     undelayed_append([H|T1],(L,ne_list),[H|T2]) if
       append(T1,L,T2).
```

```
Co-routining 2
     Conditional descriptions syntax
         <cond> ::= <variable>^(<cond>)
                    <quantified_cond>
         <quantified_cond> ::= <quantified_cond>,<quantified_cond>
                                <quantified_cond>;<quantified_cond>
                                <variable>=<cond desc>
         <cond desc> ::= <variable>
                          <type>
                         max(<type>)
                          <feat>:<cond desc>
                          <path> == <path>
                          <cond_desc>, <cond_desc>
                          <cond_desc> ;<cond_desc>
```

```
Co-routining 3
     Shared variables in conditionals
   when(X=([f]==[g]), bar(X))
         when(X=(f:Y,g:Y), bar(X))
     Narrow scope
        when(Y^{(X=(f:Y,g:Y))}, bar(X))
        foo(X) if
         Z = f:Y,
         when(Y^{(X=(f:Y,g:Y))}, bar(Y,Z))
```

```
Attaching Definite Relations: Con-
straints
     Syntax
     CondDesc *> Desc goal Goal.
     If a goal is specified in a constraint, the constraint is satisfied only if the goal
     succeeds.
     Variables occuring in the consequent are bound with scope that extends over the
      consequent and over the relation attachments.
     Example
     phrase *> (synsem:category:subcat:PhrSubcat,
                  dtr1:synsem:Synsem,
                  dtr2:synsem:category:subcat:HeadSubcat)
     goal
         append(PhrSubcat,[Synsem],HeadSubcat).
```

Attaching Structure R • Syntax	Definite Relations: Jules 1	Phrase
<rule_clause< td=""><td><pre>&gt; ::= cat&gt; <desc></desc></pre></td><td></td></rule_clause<>	<pre>&gt; ::= cat&gt; <desc></desc></pre>	
<pre>     Example     backward_ap     (synsem:     ===&gt;     cat&gt; (synsecond)     cat&gt; (synsecond)     goal&gt; app</pre>	<pre>plication rule # ., qstore:Qs) em:, qstore:Qs1), em:, qstore:Qs2), end(Qs1,Qs2,Qs).</pre>	

Attaching Definite Relations: Phrase
Structure Rules 2 Evaluation
<ul> <li>evaluated in the order they are specified</li> </ul>
<ul> <li>all possible solutions are found and the resulting instantiations are carried over to</li> <li>the rule</li> </ul>
Example
<pre>schema2 rule # (cat:(head:Head,subcat:[SubjSyn])) goal&gt; three_or_less(Comps), cats&gt; Comps, cat&gt; (cat:(head:Head,             subcat:[Subj Comps])).</pre>
<pre>three_or_less([]) if true. three_or_less([_]) if true. three_or_less([_,_]) if true. three_or_less([_,_,]) if true.</pre>

```
Attaching Definite Relations: Lexical
Rules
     Syntax:
     <lex rewrite> ::= <desc> **> <desc> if <qoal>
     passive_lex_rule ##
     (word,
      synsem:...,
      arg_st:[(loc:(cat:...,
                  cont:Cont2)), Synsem1 [List])
      * * >
     (word,
      synsem:...,
      arg_st:([Synsem1|List];Result))
     if append(([Synsem1|List]),[(loc:(cat:...,
                                     cont:Cont2))],Result)
                                     morphs
     X becomes X.
```

# Summary TRALE's definite clause extension provides a way of encoding HPSG relations

#### Summary

- TRALE's definite clause extension provides a way of encoding HPSG relations
- Various computational aspects become important in:
  - Formulating definite clauses
  - Attaching them as relational constraints

#### Summary

- TRALE's definite clause extension provides a way of encoding HPSG relations
- Various computational aspects become important in:
  - Formulating definite clauses
  - Attaching them as relational constraints
- Further reading:
  - On Prolog definite clauses: Sicstus Prolog Manual
    - Prolog Language: http://www.sics.se/sicstus/docs/latest/html/sicstus.html/Prolog-Intro.html#Prolog%20Intro
    - Predicate Index:
      - http://www.sics.se/sicstus/docs/latest/html/sicstus.html/Predicate-Index.html#Predicate%20Index

#### Summary

- On defining TRALE clauses: TRALE User Manual Chapter 5 http://www.ale.cs.toronto.edu/docs/man/ale\_trale\_man/index.html
  - Co-routining: 5.2
  - Shared variables in conditionals: 5.2.1
- On attaching goals: TRALE User Manual Chapters 4,6 Constraints: T4.3
  - Phrase structure rules: 6.4.1
  - Lexical Rules: 6.3
- On evaluating definite clause queries: Trale User Manual Appendix B, Section 2.4
- On source level debugger: Trale User Manual Appendix B, Section 2.10