The Unger Parser

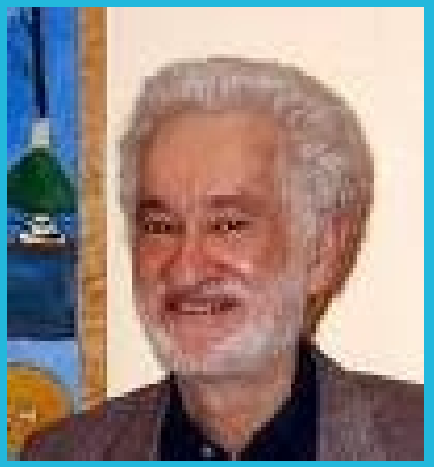
brought to you today by: Anne Brock

Outline

- Unger - the man
- Unger - the parser
- Unger's method, simple version
- some improvements
- Unger's method, including ε - rules

1. Unger: The man

Stephen H. Unger



- Politechnic Institute of Brooklyn
- doctorate at MIT
- Bell Telephone Labs
  - -research in digital systems
  - -head of development group (first electronic telephone switching system)
- since 1961: Prof. of Computer Science and Electrical Engeneering at Columbia University

-1968: the Parser.

-since: published several books.

2. The Parser

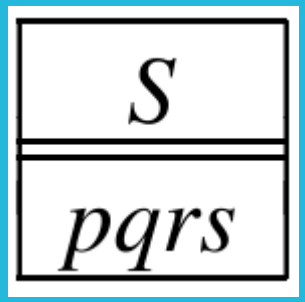
- non-directional
- top-down
- Type 2 grammars (CFG)

3. Unger's method, simplified

Input: CFG and a String/sentence, for example:

grammar: S ⟶ ABC | DE | F

'sentence': pqrs



Does S derive...

ABC | DE | F ?

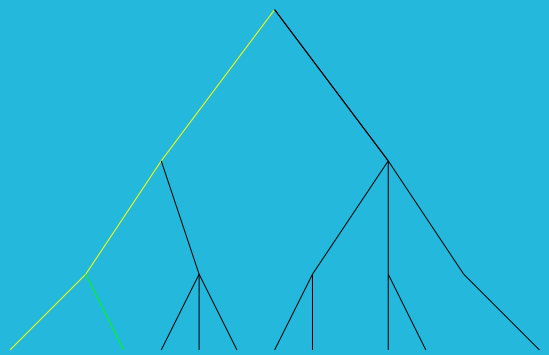
S		
A	B	C
p	q	rs
p	qr	s
pq	r	s

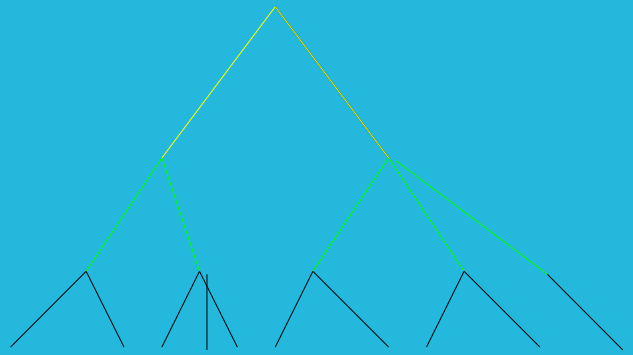
S	
D	E
p	qrs
pq	rs
pqr	s

S
F
$pqrs$

This is a search problem.

Search: depth-first or breadth-first?





A more detailed example

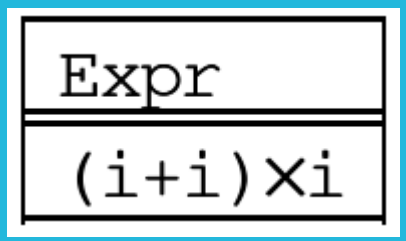
Grammar:

E -> E + T | T
T -> T x F | F
F -> (E) | i

E = Expression
T = Term
F = Factor
+, x = operators
i = operand

Input:

(i + i) x i



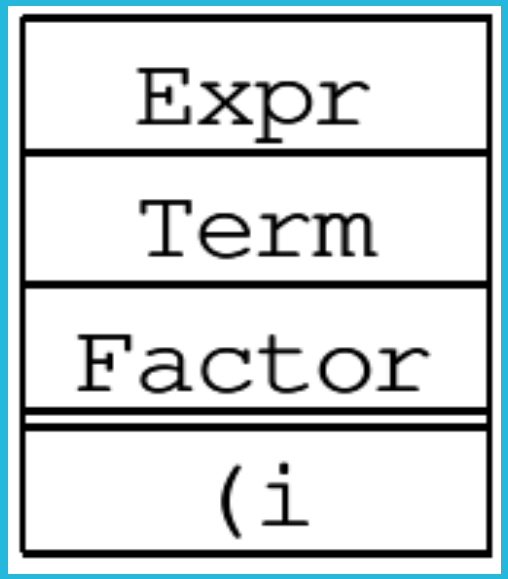
E -> E + T | T

Expr		
Expr	+	Term
(	i	+i)×i
(	i+	i)×i
(	i+i	)×i
(	i+i)	×i
(	i+i)×	i
(i	+	i)×i
(i	+i	)×i
(i	+i)	×i
(i	+i)×	i
(i+	i	)×i
(i+	i)	×i
(i+	i)×	i
(i+i	)	×i
(i+i	)×	i
(i+i)	×	i

Expr		
Term		
Term	×	Factor
(i+i)	×	i

Expr		
Expr	+	Term
(i	+	i)xi

E ->* (i ?



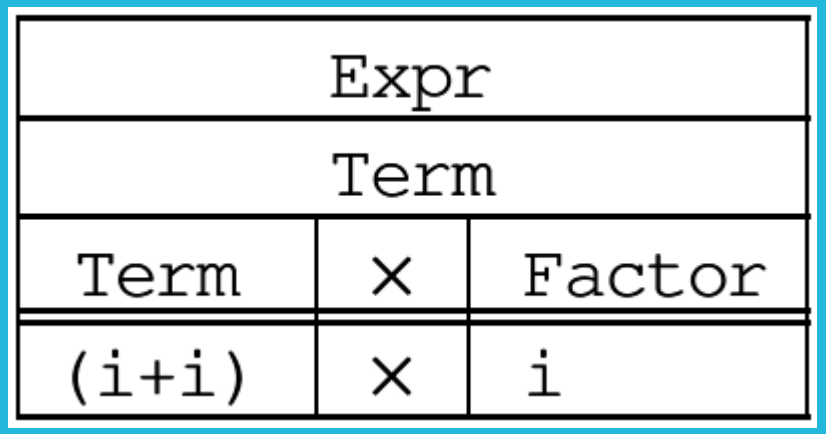
E -> E + T | T
T -> T x F | F
F -> (E) | i

fails!

to derive: (i + i) x i

E -> E + T | T

- fails!



(E -> E + T | T)
T -> T x F | F
T -> T x F | F
F -> (E) | i
F -> (E) | i

- success!

E -> E + T | T
T -> T x F | F
F -> (E) | i

```
Expr ->
Term ->
Term × Factor ->
Factor × Factor ->
( Expr ) × Factor ->
( Expr + Term ) × Factor ->
( Term + Term ) × Factor ->
( Factor + Term ) × Factor ->
( i + Term ) × Factor ->
( i + Factor ) × Factor ->
( i + i ) × Factor ->
( i + i ) × i
```

4. Room for improvement...

- consider the actual terminal symbols

- consider the length of your input

5. Unger's method with ε-rules

S -> ABC
B -> SD

try and derive:
B -> pqr

S		
A	B	C
		pqr
	p	qr
	pq	r
	pqr	
p		qr
p	q	r
p	qr	
pq		r
pq	r	
pqr		

S -> ABC
B -> SD
...

B	
S	D
	pqr
p	qr
pq	r
pqr	

What to do about it?

-> Keep a list of currently considered questions!

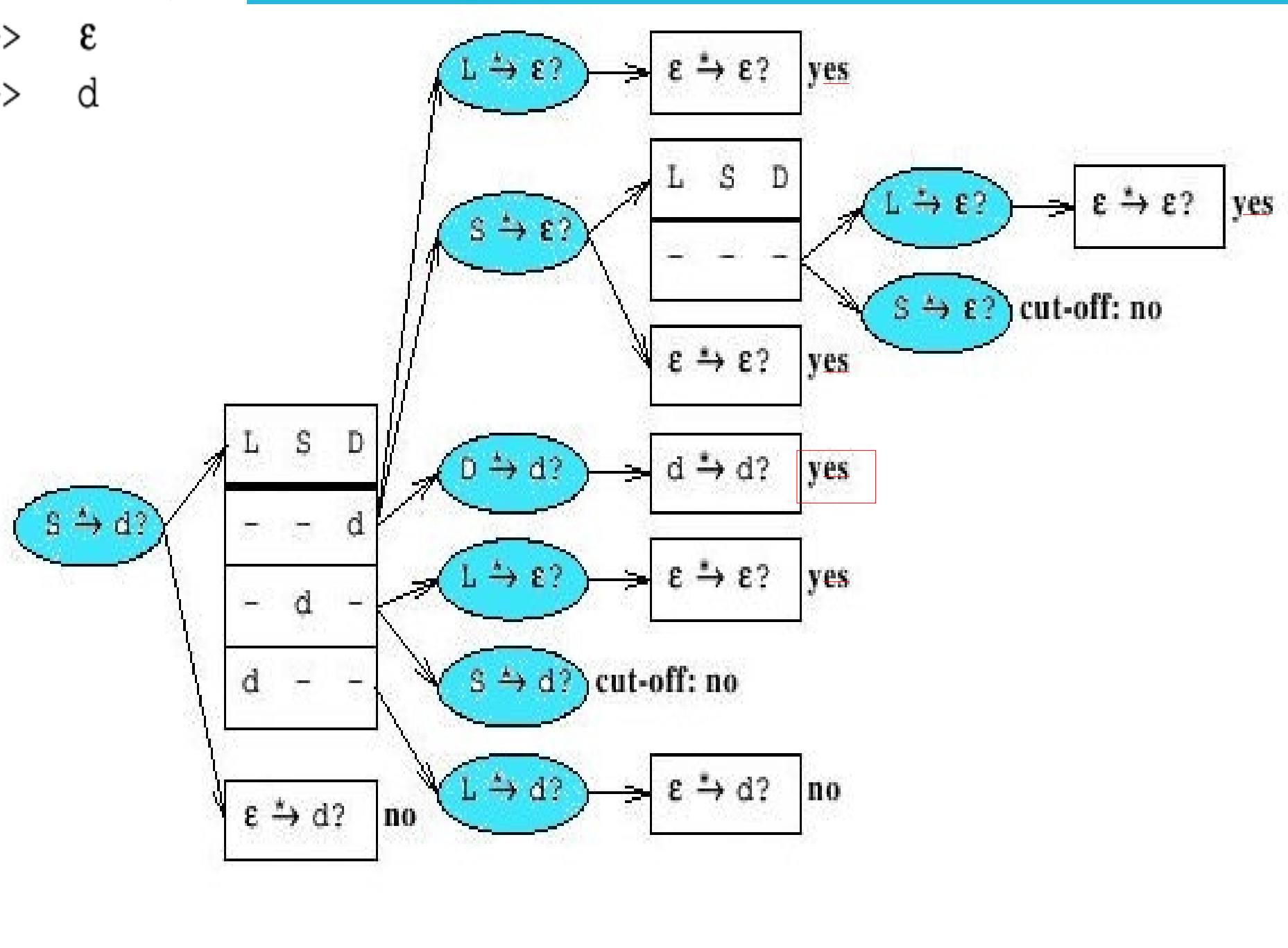
An example.

```
S   ->   LSD | ε
L   ->   ε
D   ->   d
```

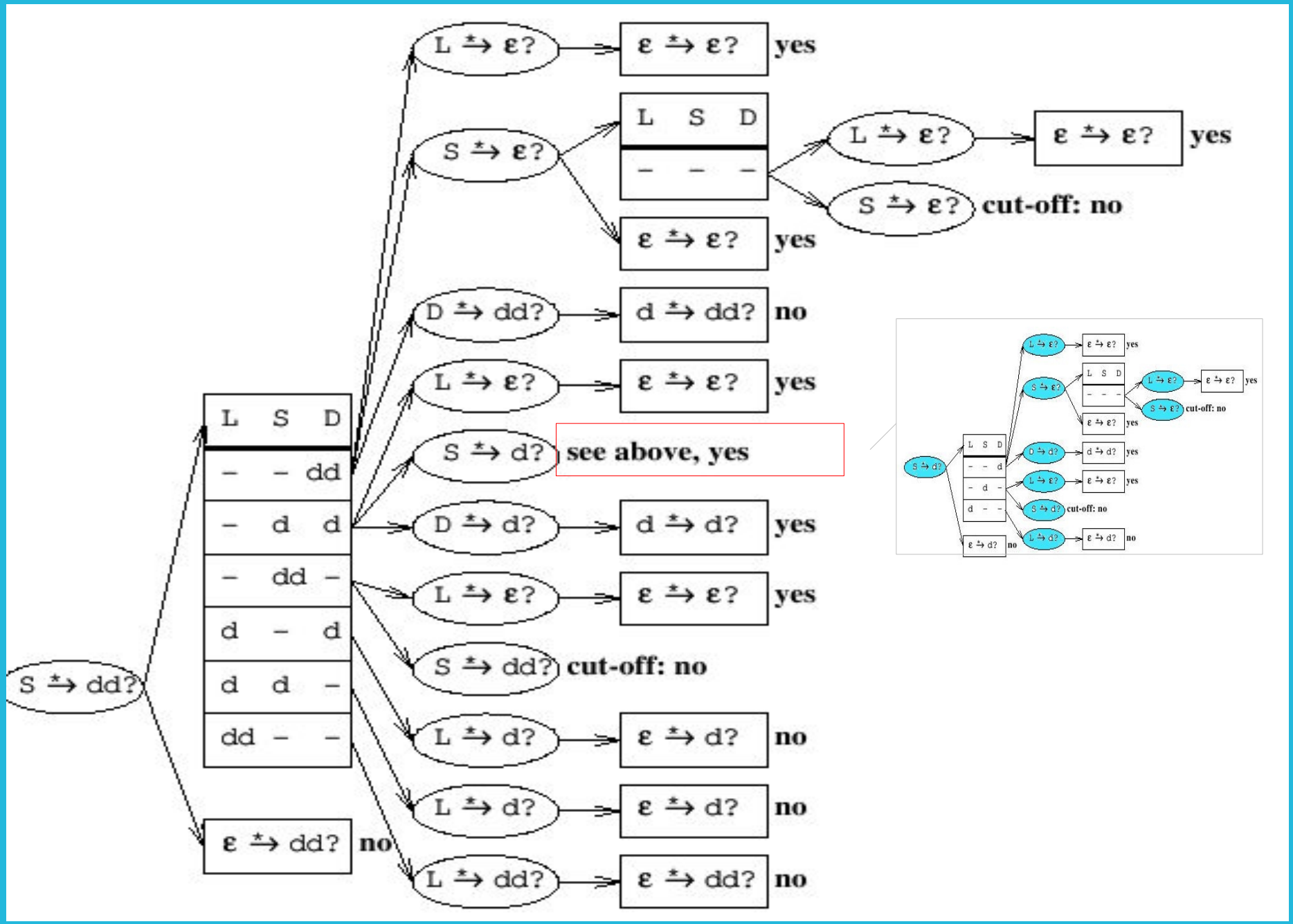
How does this grammar derive d ? dd ?

d ?

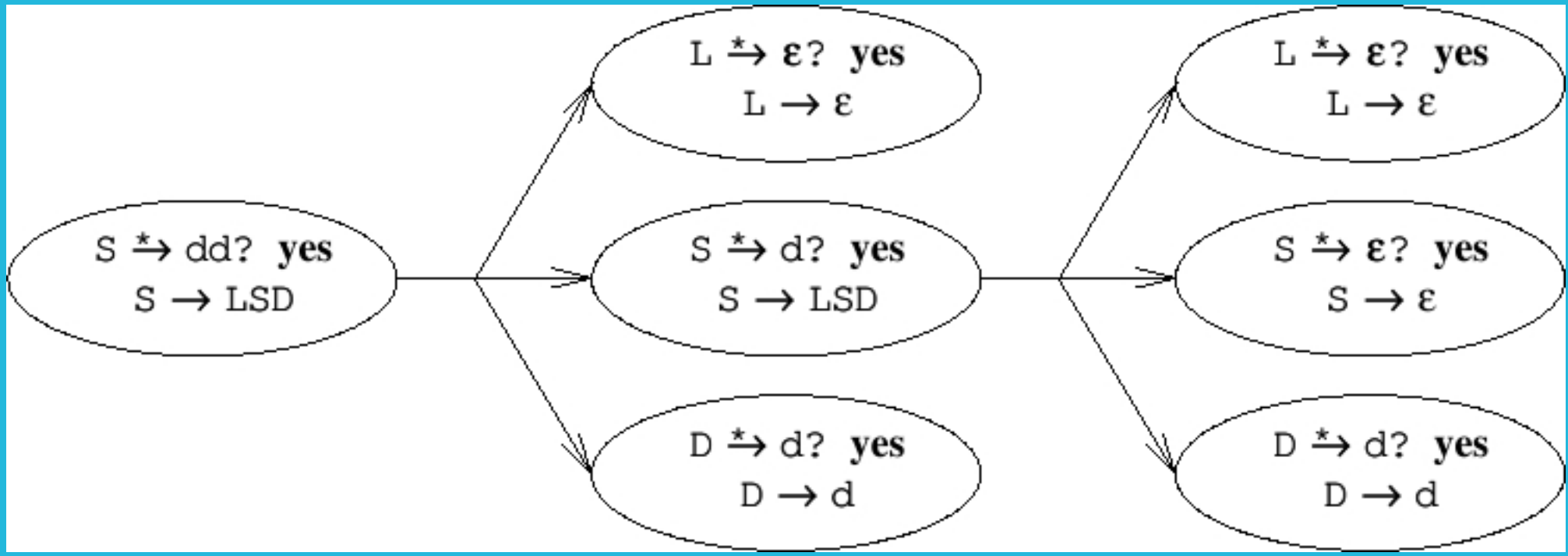
S -> LSD | ε
L -> ε
D -> d


S ⇒* d?
L S D
- - d
- d -
d - -
ε ⇒* d? no
L ⇒* ε?
ε ⇒* ε? yes
S ⇒* ε?
L S D
- - -
L ⇒* ε?
ε ⇒* ε? yes
S ⇒* ε? cut-off: no
ε ⇒* ε? yes
D ⇒* d?
d ⇒* d? yes
L ⇒* ε?
ε ⇒* ε? yes
S ⇒* d? cut-off: no
L ⇒* d?
ε ⇒* d? no

dd ?


S ⇒* dd?
L S D
– – dd
– d d
– dd –
d – d
d d –
dd – –
ε ⇒* dd? no
L ⇒* ε? ε ⇒* ε? yes
S ⇒* ε? L S D – – – yes
L ⇒* ε? ε ⇒* ε? yes
S ⇒* ε? cut-off: no
D ⇒* dd? d ⇒* dd? no
L ⇒* ε? ε ⇒* ε? yes
S ⇒* d? see above, yes
D ⇒* d? d ⇒* d? yes
L ⇒* ε? ε ⇒* ε? yes
S ⇒* dd? cut-off: no
L ⇒* d? ε ⇒* d? no
L ⇒* d? ε ⇒* d? no
L ⇒* dd? ε ⇒* dd? no

S ->* d ?



S -> LSD -> SD -> LSDD -> SDD -> DD -> dD -> dd.

Summary

The Unger parser:

- is a non-directional, top-down parser;
- will consider each possible (and impossible) solution;
- requires at least polynomial, if not exponential time;
- is slightly improved by
    - -matching input with possible derived terminals
    - -calculating possible length, special case ε
    - -remembering answers.

?

Sources

Grune, Dick and Jacobs, Ceriel 1990. *Parsing Techniques. A Practical Guide.* New York: Ellis Horwood Limited.

Lukasz Kwiatowski. *Reconciling Unger's parser as a top-down parser for CF grammars for experimental purposes.* http://www.cs.vu.nl/~steven/

pictures from:

www.cs.columbia.edu/async/images/unger.jpg

http://pinker.wjh.harvard.edu/photos/cambridge_boston/pages/trees%20in%20Cambridge%20Common.htm